

# Chapter 17

17.1 Let there be three items A, B and C. The raw materials required for these products are  $r_1$  and  $r_2$ . The daily allotment of raw materials is 90 and 80 respectively. How the raw materials should be allotted so that the profit is maximized.

Product	Raw material		Profit per unit
	$r_1$	$r_2$	
A	2	2	140
B	3	5	120
C	5	5	250
Daily Allotment	10	14	

**Solution:**

Based on example 17.1, one can formulate LPP as

$$3x_1 + 3x_3 + 5x_3 \leq 10$$

$$2x_1 + 5x_2 + 5x_3 \leq 14$$

$$x_1, x_2, x_3 \geq 0$$

and the objective function is maximize

$$Z = 140x_1 + 120x_2 + 250x_3$$

17.2 There are three materials N1, N2 and N3. The need of a patient is at least 14 and 18. How to mix the following nutrients for the given cost consideration?

Product	Raw material			Profit per unit
	$N_1$	$N_2$	$N_3$	
$F_1$	6	2	4	9
$F_2$	2	4	6	15
$F_3$	5	3	7	23
Need	14	18	24	—

$$6x_1 + 2x_3 + 5x_3 \leq 14$$

$$2x_1 + 4x_2 + 3x_3 \leq 18$$

$$4x_1 + 6x_2 + 7x_3 \leq 24$$

$$x_1, x_2, x_3 \geq 0$$

and the objective condition is

$$Z = 9x_1 + 15x_2 + 23x_3$$

17.3 Formulate an LPP for finding the shortest path in a given graph.

**Solution:** The input graph G is taken with source s and goal t. The edge weight represents the path from s to v.

The LPP formulation would be as follows: Let  $X_i$  be the distance between source and vertex  $i$ , then the LPP would be like

$$\text{Minimize } X_t$$

Subject to the conditions  $X_i$ .

17.4 Formulate an LPP for finding the minimum cost spanning tree in a graph.

**Solution:** The objective is to minimize the cost.

minimize  $\sum_{i,j \in E} c_{ij}x_{ij}$ , Here  $c$  is the cost linking two vertices liked by edge  $(i,j)$ . One constraint

is, there are  $N-1$  edges in  $T$ . The second constraint is any subset  $k$  involves only  $k-1$  edges so that there is no cycle.

This is the structure of LPP problem.

17.5 Solve the following linear programming using the graphical method.

$$\text{Maximize } Z = 2x_1 + 4x_3$$

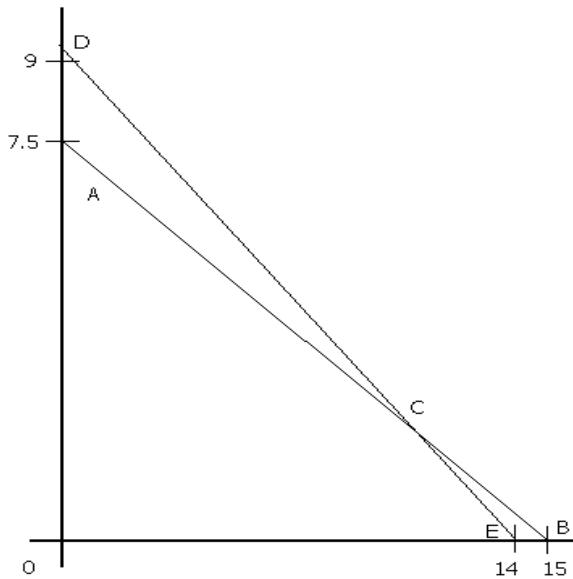
Subjected to

$$x_1 + 2x_2 \leq 15$$

$$2x_1 + 3x_2 \leq 28$$

$$x_1, x_2 \geq 0$$

**Solution:**



Point	X	Y	Value of Z
<b>0</b>	0	0	0
<b>A</b>	0	7.5	30
<b>B</b>	15	0	30
<b>C</b>	11	2	30
<b>D</b>	0	9.33	37.33
<b>E</b>	14	0	23

17.6 Formulate a LPP model for maximum bipartite matching problem.

**Solution:** Here, the vertices  $V$  is divided into two sets with edge connecting between vertices of two sets. The LPP formulation would be as follows:

One variable for every edge. For example, if vertex 1 is linked with A, then a variable is created as  $A_1$ . One equality condition is created per vertex. The goal is to find the maximum cardinality matching.

So the problem shape is Maximize {Variables connecting Edges}

Subjected to All edges connected with vertex added with equality one.

17.7 Solve the following LPP using graphical method.

$$\text{Maximize } Z = 24x_1 + 14x_2$$

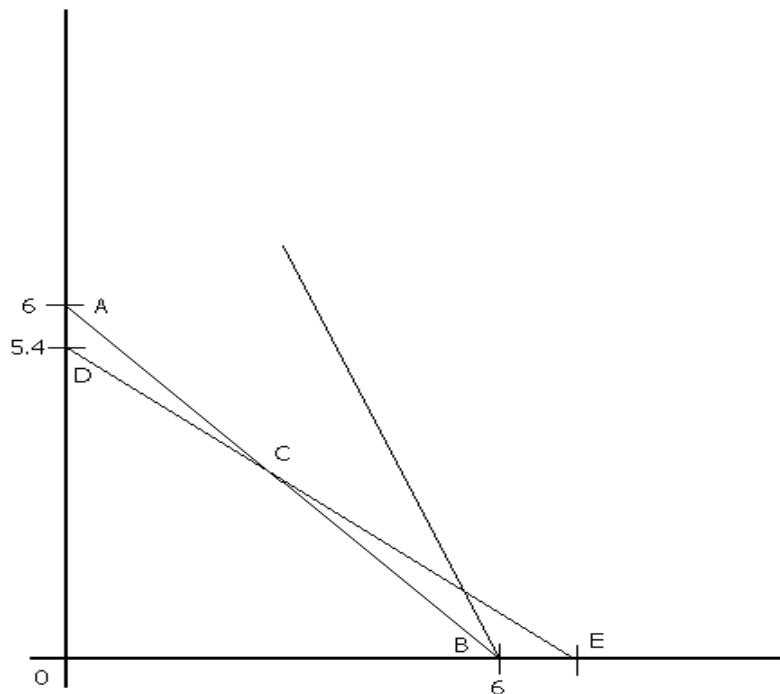
Subjected to

$$8x_1 + 8x_2 \leq 48$$

$$9x_1 + 12x_2 \leq 64$$

$$x_1, x_2 \geq 0.$$

**Solution:**



	X	Y	Value of Z
<b>0</b>	0	0	0
<b>A</b>	0	6	8.40
<b>B</b>	6	0	14.40
<b>C</b>	2.66	3.33	110.66
<b>D</b>	0	5.33	74.66
<b>E</b>	7.11	0	170.66

$\therefore$  Answer is  $x_1 = 7.1, x_2 = 0$

17.8 Solve the following LPP using graphical method.

Minimize  $Z = 12x_1 + 8x_2$

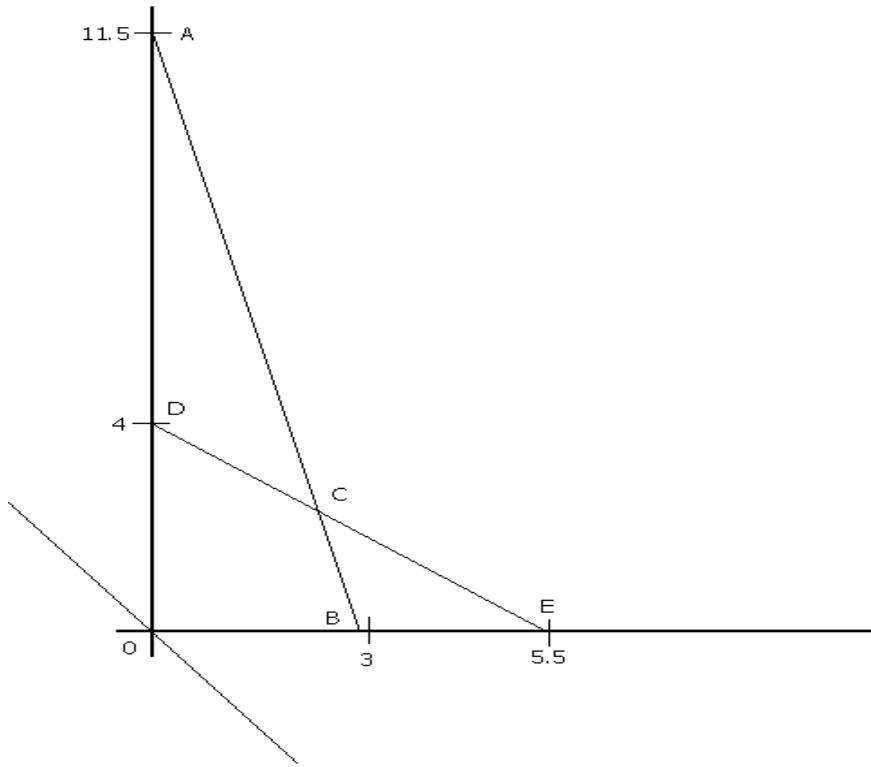
Subjected to

$$8x_1 + 2x_2 \leq 23$$

$$9x_1 + 12x_2 \leq 48$$

$$x_1, x_2 \geq 0$$

*Solution:*



Point	X	Y	Value of Z
0	0	0	0
A	0	11.5	92.00
B	2.87	0	34.50
C	2.31	2.27	45.85
D	0	4	32.00
E	5.5	0	64.00

$\therefore$  Answer is  $x = 0, y = 4$ .

17.9 Solve the following LPP using the simplex method.

$$\text{Minimize } Z = 4x_1 + 3x_2 + 8x_3$$

$$\text{Subjected to } x_1 + 6x_2 + 12x_3 \leq 25$$

$$2x_1 + 8x_2 + 16x_3 \leq 42$$

$$x_1, x_2, x_3 \geq 0$$

**Solution:**

$$4x_1 + 3x_2 + 8x_3 + 0x_4 + 0x_5$$

Subjected to

$$x_1 + 6x_2 + 12x_3 + x_4 = 25$$

$$2x_1 + 8x_2 + 16x_3 + x_5 = 42$$

$$x_1, x_2, x_3, x_4, x_5 \geq 0$$

Table			4	3	8	0	0
Base	$C_b$	$p_0$	$p_1$	$p_2$	$p_3$	$p_4$	$p_5$
$p_4$	0	25	1	6	12	1	0
$P_5$	0	42	3	8	16	0	1
$Z$	0	-4	-3	-8	0	0	0

The leaving variable is  $p_4$  and entering variable is  $p_3$

Table			4	3	8	0	0
Base	$C_b$	$p_0$	$p_1$	$p_2$	$p_3$	$p_4$	$p_5$
$P_3$	8	$\frac{25}{12}$	$\frac{1}{12}$	$\frac{1}{2}$	1	$\frac{1}{12}$	0
$P_5$	0	$\frac{26}{3}$	$\frac{2}{3}$	0	0	$\frac{4}{3}$	1
$Z$	0	$\frac{-50}{3}$	$\frac{-10}{3}$	1	0	$\frac{2}{3}$	0

The leaving variable is  $p_5$  and entering variable is  $p_1$

<b>Table 3</b>			<b>4</b>	<b>3</b>	<b>8</b>	<b>0</b>	<b>0</b>
<b>Base</b>	$C_b$	$p_0$	$p_1$	$p_2$	$p_3$	$p_4$	$p_5$
$P_3$	8	1	0	$\frac{1}{2}$	1	$\boxed{\frac{1}{4}}$	$-\frac{1}{8}$
$P_1$	4	13	1	0	0	-2	$\frac{3}{2}$
$Z$	0	60	0	1	0	-6	5

The leaving variable is  $p_3$  and entering variable is  $p_4$

<b>Table 4</b>			<b>4</b>	<b>3</b>	<b>8</b>	<b>0</b>	<b>0</b>
<b>Base</b>	$C_b$	$p_0$	$p_1$	$p_2$	$p_3$	$p_4$	$p_5$
$P_4$	0	4	0	2	4	1	$-\frac{1}{2}$
$P_1$	4	21	1	4	8	0	$\frac{1}{2}$
$Z$	0	84	0	13	24	0	2

$\therefore$  The Optimal solution is  $z = 84$

$$x_1 = 21$$

$$x_2 = 0$$

$$x_3 = 0$$

17.10 Solve the following LPP using simplex method

$$\text{Maximize } Z = 12x_1 + 24x_2 + 23x_3 + x_4$$

$$\text{Subjected To } 2x_1 + 4x_2 + 5x_3 + 6x_4 \leq 43$$

$$3x_1 + 4x_2 + 5x_3 + 5x_4 \leq 64$$

$$x_1, x_2, x_3, x_4 \geq 0$$

**Solution:**

$$12x_1 + 24x_2 + 23x_3 + x_4 + 0x_5 + x_6$$

$$2x_1 + 4x_2 + 5x_3 + 6x_4 + x_5 = 43$$

$$3x_1 + 4x_2 + 5x_3 + 5x_4 + x_6 = 64$$

$$x_1, x_2, x_3, x_4 \geq 0$$

Table 1			12	24	23	1	0	0
Base	$C_b$	$p_0$	$p_1$	$p_2$	$p_3$	$p_4$	$p_5$	$P_6$
$P_5$	0	43	2	4	5	6	1	0
$P_6$	0	64	3	4	5	5	0	1
Z	0	0	-12	-21	-23	-1	0	0

The leading is  $p_5$  and entering variable  $p_2$ .

Table 2			12	24	23	1	0	0
Base	$C_b$	$p_0$	$p_1$	$p_2$	$p_3$	$p_4$	$p_5$	$P_6$
$P_2$	24	10.25	0.5	1	1.25	1.5	0.25	0
$P_6$	0	21	1	0	0	-1	-1	1
Z	0	258	0	0	7	35	6	0

There is infinitely many values of  $x_1, x_2, x_3, x_4$  for the optimal value of  $Z = 258$ , which are contained in the region of space  $12x_1 + 24x_2 + 23x_3 + x_4 = 258$  that satisfies the constraints.

$$\text{Answer is } x_1 = 0$$

$$x_2 = 10.75$$

$$x_3 = 0$$

$$x_4 = 0$$

17.11 Solve the following LPP using the simplex method

$$Z = 8x_1 + 9x_2 + 2x_3$$

Subjected To

$$3x_1 + 4x_2 + 5x_3 \leq 16$$

$$2x_1 + 3x_2 + x_3 \leq 24$$

$$x_1, x_2, x_3 \geq 0$$

**Solution:**

$$\text{Maximize } 8x_1 + 9x_2 + 2x_3 + x_4 + x_5$$

Subjected to

$$3x_1 + 4x_2 + 5x_3 + x_4 = 16$$

$$2x_1 + 3x_2 + x_3 + x_5 = 24$$

$$x_1, x_2, x_3, x_4, x_5 \geq 0$$

<b>Table 1</b>			<b>8</b>	<b>9</b>	<b>2</b>	<b>0</b>	<b>0</b>
<b>Base</b>	$C_b$	$p_0$	$p_1$	$p_2$	$p_3$	$p_4$	$p_5$
$P_4$	0	16	3	4	5	1	0
$P_5$	0	24	2	3	1	0	1
$Z$		0	-5	-9	-2	0	0

The leaving variable is  $p_4$  and entering variable is  $p_2$ .

<b>Table 2</b>			<b>8</b>	<b>9</b>	<b>2</b>	<b>0</b>	<b>0</b>
<b>Base</b>	$C_b$	$p_0$	$p_1$	$p_2$	$p_3$	$p_4$	$p_5$
$P_2$	9	4	$\frac{3}{4}$	1	$\frac{5}{4}$	$\frac{1}{4}$	0
$P_5$	0	12	$-\frac{1}{4}$	0	$-\frac{11}{4}$	$-\frac{3}{4}$	1
$Z$		36	$-\frac{5}{4}$	0	$\frac{37}{4}$	$\frac{9}{4}$	0

The leaving variable is  $p_2$  and entering variable is  $p_1$ .

<b>Table 3</b>			<b>8</b>	<b>9</b>	<b>2</b>	<b>0</b>	<b>0</b>
<b>Base</b>	$C_b$	$p_0$	$p_1$	$p_2$	$p_3$	$p_4$	$p_5$
$P_4$	8	$\frac{16}{3}$	1	$\frac{4}{3}$	$\frac{5}{3}$	$\frac{1}{3}$	0

<b><math>P_5</math></b>	0	$\frac{40}{3}$	0	$\frac{1}{3}$	$\frac{-7}{3}$	$\frac{-2}{3}$	1
<b>Z</b>	0	$\frac{128}{3}$	0	$\frac{5}{3}$	$\frac{34}{3}$	$\frac{8}{3}$	0

∴ The optimal solution is

$$Z = \frac{128}{3}$$

$$x_1 = \frac{16}{3}$$

$$x_2 = 0$$

$$x_3 = 0$$

17.12 Solve the following LPP using the two-table simplex method.

$$\text{Minimize } Z = 2x_1 + 4x_2 + 3x_3$$

Subjected to

$$x_1 + x_2 + x_3 \leq 18$$

$$x_1 + 3x_2 + 2x_3 \leq 28$$

$$x_1, x_2, x_3 \geq 0$$

**Solution:**

$$\text{Maximize } -2x_1 - 4x_2 - 3x_3 + x_4 + x_5 + x_6$$

Subjected to

$$x_1 + x_2 + x_3 + x_4 = 18$$

$$x_1 + 3x_2 + 2x_3 + x_5 = 28$$

$$x_1, x_2, x_3, x_4, x_5, x_6 \geq 0$$

Construct phase 1

<b>Table 1</b>			<b>0</b>	<b>0</b>	<b>0</b>	<b>0</b>	<b>-1</b>	<b>-1</b>
<b>Base</b>	$C_b$	$p_0$	$p_1$	$p_2$	$p_3$	$p_4$	$p_5$	$P_6$
$P_6$	-1	18	1	1	1	-1	0	1
$P_5$	-1	28	1	3	2	0	1	0
<b>Z</b>	0	-46	-2	-4	-3	1	0	0

The leaving variable is  $p_5$  and entering variable is  $p_2$ .

<b>Table 2</b>			<b>0</b>	<b>0</b>	<b>0</b>	<b>0</b>	<b>-1</b>	<b>-1</b>
<b>Base</b>	$C_b$	$p_0$	$p_1$	$p_2$	$p_3$	$p_4$	$p_5$	$P_6$
$P_6$	-1	$\frac{26}{3}$	$\frac{2}{3}$	0	$\frac{1}{3}$	-1	$-\frac{1}{3}$	1
$P_2$	0	$\frac{25}{3}$	$\frac{1}{3}$	1	$\frac{2}{3}$	0	$\frac{1}{3}$	0
$Z$	0	$-\frac{26}{3}$	$-\frac{2}{3}$	0	$-\frac{1}{3}$	1	$\frac{4}{3}$	0

The leaving variable is  $p_6$  and entering variable is  $p_1$ .

<b>Table 3</b>			<b>0</b>	<b>0</b>	<b>0</b>	<b>0</b>	<b>-1</b>	<b>-1</b>
<b>Base</b>	$C_b$	$p_0$	$p_1$	$p_2$	$p_3$	$p_4$	$p_5$	$P_6$
$P_1$	0	13	1	0	$\frac{1}{2}$	$-\frac{3}{2}$	$-\frac{1}{2}$	$\frac{3}{2}$
$P_2$	0	5	0	1	$\frac{1}{2}$	$\frac{1}{2}$	$\frac{1}{2}$	$-\frac{1}{2}$
$Z$		0	0	0	0	0	1	1

There may be any solution, so one can continue with phase II to calculate it.

<b>Table 1</b>			<b>-2</b>	<b>-4</b>	<b>-3</b>	<b>0</b>
<b>Base</b>	$C_b$	$p_0$	$p_1$	$p_2$	$p_3$	$p_4$
$P_1$	-2	13	1	0	$\frac{1}{2}$	$-\frac{3}{2}$
$P_2$	-4	5	0	1	$\frac{1}{2}$	$\frac{1}{2}$
$Z$		-46	0	0	0	1

There are many solutions. One solution is

$$x_1 = 13$$

$$x_2 = 5$$

$$x_3 = 0$$

17.13 Find the dual at the following LPP and verify it

$$\text{Maximize } Z = 2x_1 + 4x_2$$

Subjected to

$$x_1 + 2x_2 \leq 15$$

$$2x_1 + 3x_2 \leq 28$$

$$x_1, x_2 \geq 0$$

**Solution:**

The dual problem can be given as

$$\text{Minimize } Z' = 15y_1 + 28y_2$$

Subjected to the condition

$$y_1 + 2y_2 \leq 2$$

$$2y_1 + 3y_2 \leq 4$$

$$y_1, y_2 \geq 0.$$

The verification part is, the dual of the dual is its primal. It is the proof.

17.14 Find the dual of the following LPP and verify it.

$$\text{Maximize } Z = 12x_1 + 8x_2$$

Subjected to

$$8x_1 + 2x_2 \leq 23$$

$$9x_1 + 12x_2 \leq 48$$

$$x_1, x_2 \geq 0$$

**Solution:**

The dual problem can be written as

$$\text{Minimize } Z = 23y_1 + 48y_2$$

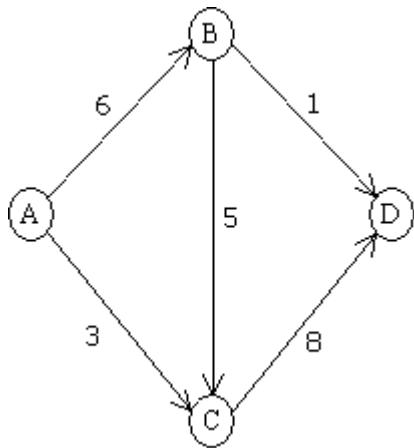
$$8y_1 + 9y_2 \leq 12$$

$$2y_1 + 12y_2 \leq 8$$

$$y_1, y_2 \geq 0$$

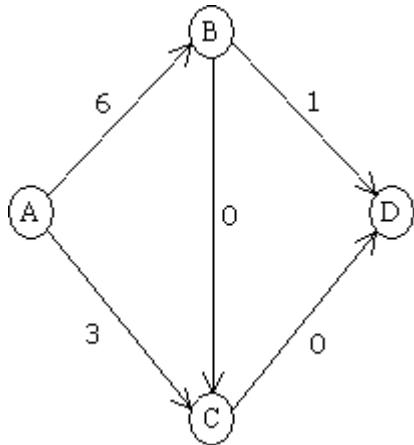
Hence  $y_1, y_2$  is unrestricted with respect to sign.

17.15 Solve the following problem using Ford–Fulkerson algorithm.



**Solution:**

The initial flow graph is given as follows.



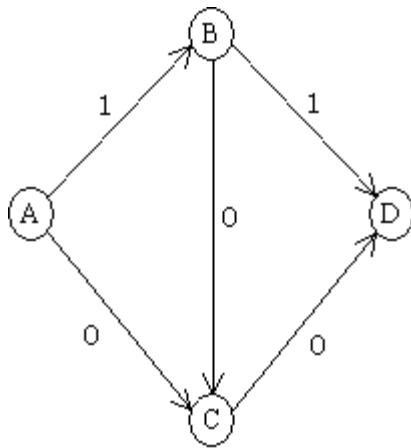
The commodity can be increased by 3.

Construct one augment path  $A \rightarrow B \rightarrow D$ .

Edges	Total Capacity	Current Load	Excess Capacity
$A \rightarrow B$	6	0	6
$B \rightarrow D$	0	1	1

Minimum { 1, 6 } = 1

$\therefore$  Increase the load by 1.

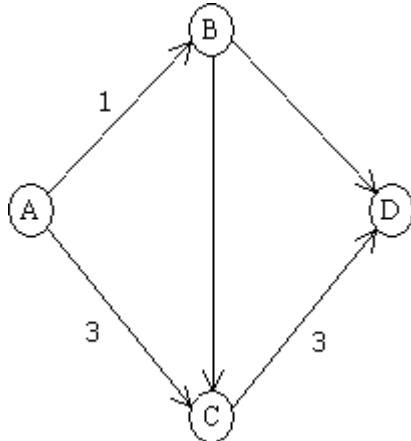


Next augment path is ACD

Edges	Total Capacity	Current Load
$A \rightarrow C$	3	0
$C \rightarrow D$	8	0

Minimum { 3, 8 } = 3

∴ The updated graph would be

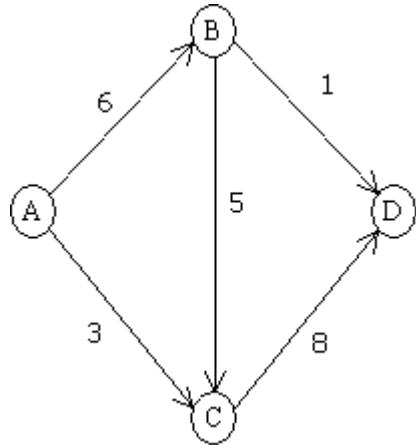


The next augment path is ABCD.

Edges	Total Capacity	Current Load	Excess Capacity
$A \rightarrow B$	6	1	5
$B \rightarrow C$	5	0	5

<b>C → D</b>	8	3	5
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Minimum { 5, 5, 5 }

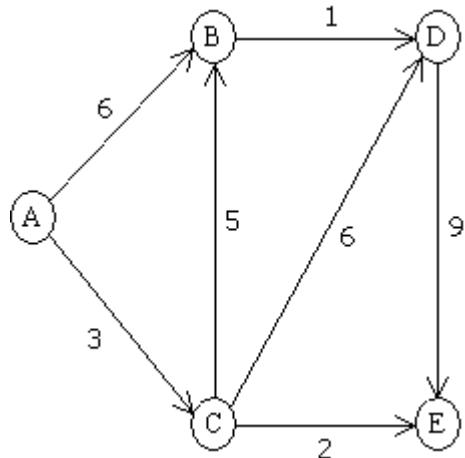


Incidentally the maximum flow is same on the original graph.

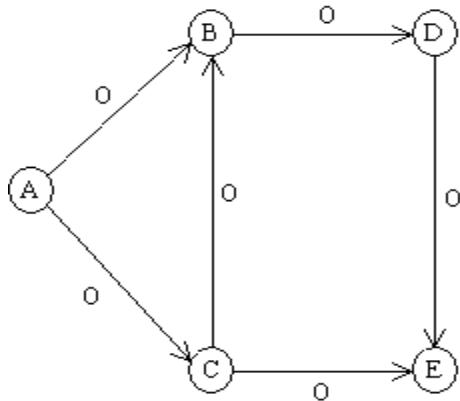
17.16 Solve the following problem.

**Solution:**

The initial graph is



The initial augment path is



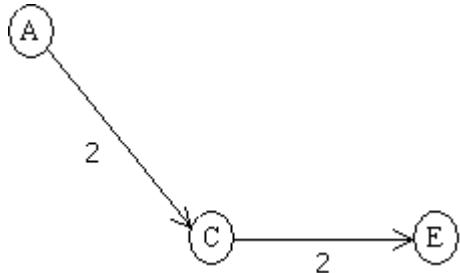
A C E

The minimum is completed as  $A \rightarrow C$

Edges	Total Capacity	Current Load
$A \rightarrow C$	3	0
$C \rightarrow D$	2	0

Minimum  $\{ 2, 3 \} = 2$ .

$\therefore$  Increase capacity by 2 to get this graph.

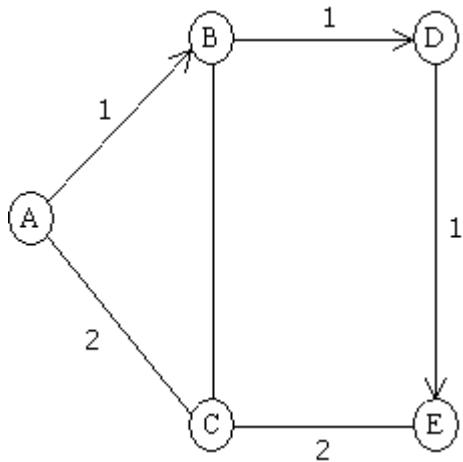


The next augment path is A B D E

Edges	Total Capacity	Current Load	Excess Capacity
$A \rightarrow B$	6	0	6
$B \rightarrow D$	1	0	1
$B \rightarrow E$	9	0	9

Minimum  $\{ 6, 1, 9 \} = 1$ .

∴ Increase by 1.

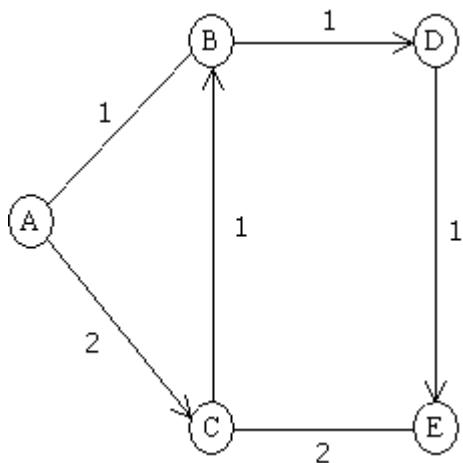


The next augment path is A C B D E

Edges	Total Capacity	Current Load	Excess Capacity
A → C	3	2	1
C → B	5	0	5
B → D	1	1	0
D → E	9	1	8

Minimum is {0}

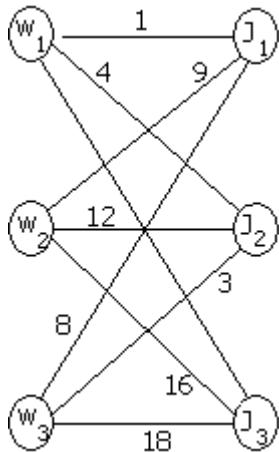
The final graph is



The minimum flow is 3.

17.17 Let there be 3 workers and three jobs, the cost matrix for which is as follows:

Workers \ Jobs	1	2	3
1	2	4	8
2	9	12	16
3	1	3	18



One can apply Hungarian method now.

17.18 Solve the stable marriage problem for 3 boys and 3 girls with the following pref. matrix.

<b><math>B_0</math></b>	<b><math>G_0</math></b>	<b><math>G_2</math></b>	<b><math>G_1</math></b>
<b><math>B_1</math></b>	$G_2$	$G_0$	$G_1$
<b><math>B_2</math></b>	$G_2$	$G_1$	$G_0$
<b><math>G_0</math></b>	$B_1$	$B_0$	$B_2$
<b><math>G_1</math></b>	$B_2$	$B_1$	$B_0$
<b><math>G_2</math></b>	$B_1$	$B_0$	$B_2$

Show the stable matching pairs.

***Solution:***

## Round I

- $B_0$  prefers  $G_0$
- $G_0$  is unmatched
- $B_0$  becomes pair with  $G_0$
- $B_1$  prefers  $G_2$
- $G_2$  previously unmatched
- $B_1$  pair with  $G_2$
- $B_2$  prefers  $G_2$
- $G_2$  is already paired with  $B_1$
- $G_2$  rank  $B_1$  as 0
- $G_2$  rank  $B_2$  as 2
- $B_2$  moves on to  $G_1$

## RoundII

- $B_0$  prefers  $G_0$
- $G_0$  already matched to  $B_0$
- $B_1$  prefers  $G_2$
- $G_2$  is already matched to  $B_1$
- $B_2$  prefers  $G_1$
- $G_1$  is previously unmatched
- $B_2$  becomes pair with  $G_1$

$\therefore$  Result:

$B_0$  with  $G_0$

$B_1$  with  $G_2$

$B_2$  with  $G_1$